Week 2 Wednesday Review Quiz

Q1 Iterated transition function 1 Point

For some proofs it is useful to define the iterated transition function of a finite automaton. Given a finite automaton $(Q, \Sigma, \delta, q_0, F)$ the iterated transition function δ^* has domain $Q \times \Sigma^*$ and codomain Q and, on input (q, w), outputs the state that the machine is in when starting in state q and processing the entire string w character-by-character according to the transition function δ .

Since the string w is built up recursively, we will define the iterated transition function recursively. Choose the correct option to fill in this definition:

The basis case is for w=arepsilon and we define

$$\delta^*((q,arepsilon)) \ \odot arepsilon \ arep$$

=

The recursive step is when w=ua for $u\in \Sigma^*$ and $a\in \Sigma$, and we define

$$egin{aligned} \delta^*((q,w)) &= \ &\bigcirc \delta((\delta^*((q,u)),a)) \ &\bigcirc \delta((\delta((q,u)),a)) \end{aligned}$$

Q2 Iterated transition function example 2 Points

Let's consider the finite automaton whose state diagram is below. Notice that we can deduce the alphabet from the diagram: the arrow labels are the symbols in the alphabet so $\Sigma = \{a, b\}$. Let's refer to the transition function of this automaton by δ .



Q2.2 (c) 1 Point What is $\delta^*((q0, a))$? $\bigcirc q0$ $\bigcirc q1$ $\bigcirc q2$ $\bigcirc a$ $\bigcirc b$ \bigcirc Undefined

What is $\delta^*((q0,abb))$?

 $\bigcirc q0$

 $\bigcirc q1$

 \bigcirc q2

 $\bigcirc a$

 $\bigcirc b$

 \bigcirc Undefined

Q3 General constructions

3 Points

Fix the alphabet $\Sigma = \{a, b\}$. For each positive integer n, define L_n to be the language over Σ given by $L_n = \{w \in \Sigma^* \mid |w| \text{ is an integer multiple of } n\}$

Q3.1 Example 1 Point

Select all and only the languages below in which abbaab is an element.

$\Box L_1$
$\Box L_2$
$\Box L_3$
$\Box L_4$
$\Box L_5$
$\Box L_6$
$\Box L_7$
$\Box L_8$
$\Box L_9$
$\Box L_{10}$
$\Box L_{11}$
$\Box L_{12}$

Q3.2 Automata and \$\$L_n\$\$ 1 Point

True or false: for each n, there is some finite automaton that recognizes L_n .

 \bigcirc True

 \bigcirc False



Q3.3 A related language 1 Point

True or false: There is a finite automaton that recognizes the set of all strings over $\{a, b\}$ with odd length.

 \bigcirc True

 \bigcirc False



Q4 DFA construction 2 Points

Consider an arbitrary finite automaton $M = (Q, \{a, b\}, \delta, q_0, F)$ and let's call the language recognized by this finite automaton L.

We can define a new finite automaton which recognizes the collection of strings that result from taking each string in L and replacing each a in the string with 0 and each b in the string with 1. For example, if $L = \{a, aab\}$, then this process would produce the new language $\{0, 001\}$.

Informally: the construction is to keep the states and arrows more or less the same, but change the labels so that the label a on an arrow is replaced by the label 0 and the label b on an arrow is replaced by the label 1.

Fill in the formal definition below;

The new machine is $M' = (Q', \Sigma', \delta', q', F')$ where

 $egin{aligned} Q' &= & & \ & \bigcirc Q & & \ & \bigcirc \overline{Q}, ext{ aka } Q^c, ext{ aka the complement of } Q & & \ & \bigcirc Q imes Q & & \ & \searrow' = & & \ & \bigcirc \{a, b\} & & \ & \bigcirc \{0, 1\} & & \ & \bigcirc \{0, 1, a, b\} & & \end{aligned}$

 $\delta':Q' imes\Sigma' o Q'$ is defined by $\delta'((q,0))=\delta((q,a))$ and $\delta'((q,1))=\delta((q,b))$ for each $q\in Q.$

q' =
\bigcirc 0
\bigcirc 1
$\bigcirc a$
$\bigcirc b$
$\bigcirc q$
$\bigcirc q_0$
F' =
\bigcirc F
$igcap \overline{F}$, aka F^c , aka the complement of F
\bigcirc $F imes F$



Q5 Closure 2 Points

A set X is said to be **closed** under an operation OP if, for any elements in X, applying OP to them gives an element in X. For example, the set of integers is closed under multiplication because if we take any two integers, their product is also an integer.

For each of the sentences below, (1) first determine if it is a closure claim, and, if it is, then (2) determine if the sentence is true or false.

Concatenating two strings over the alphabet Σ gives a string over the alphabet Σ

- \bigcirc Not a closure claim.
- \bigcirc Is a closure claim, but false.
- \bigcirc Is a closure claim, and true.

The intersection of two infinite sets of integers is an infinite set of integers.

- \bigcirc Not a closure claim.
- \bigcirc Is a closure claim, but false.
- \bigcirc Is a closure claim, and true.

Save Answer

Q6 Feedback 0 Points

Any feedback or questions about today's material or comments you'd like to share? (Optional; not for credit)



Week 2 Friday Review Quiz

Student Name

Search students by name or email...

Q1 Strings accepted by NFA

2 Points

Select all (and only) the strings of length 3 that are accepted by the NFA over the alphabet $\{0,1\}$ with state diagram:

•



Q2 Strings accepted by NFA 2 Points

Select all (and only) the strings of length 3 that are accepted by the NFA over the alphabet $\{0,1\}$ with state diagram:

1 0		
Q0	1	 →q1
000		
001		
010		
011		
☐ 100		
101		
□ 110		
□ 111		

Q3 NFA and regular expressions 6 Points

Consider the NFA over the alphabet $\{a,b\}$ with state diagram:



Q3.1 Language recognized by NFA 3 Points

Pick the regular expression that describes the language of this NFA.

- $\bigcirc (a \cup b)^* ab$
- $\bigcirc (a \cup b)^* ab(a \cup b)^*$
- $\bigcirc (a \cup b)^* ab(a \cup b)^* b$
- $\bigcirc (a \cup b)^*a(a \cup b)^*b$
- $\bigcirc (a \cup b)^* a(a \cup b)b$
- \bigcirc None of the above

Q3.2 Languages described by regular expressions 3 Points

Select all and only regular expressions that describe a subset of the language recognized by the NFA.



Save Answer

Q4 Feedback 0 Points

Any feedback or questions about today's material or comments you'd like to share? (Optional; not for credit)



Save All Answers

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